

The Lattice Structure of Multiplicative Congruential Pseudo-Random Vectors*

By W. A. Beyer, R. B. Roof and Dorothy Williamson

Abstract. The lattice structure of points in an n -dimensional space produced by an appropriate grouping of pseudo-random numbers obtained from multiplicative congruential generators is discussed. Examples are given for $2 \leq n \leq 6$. The work is based on the theory of the reduction of positive quadratic forms in n variables.

1. Introduction. There have recently appeared several articles [3], [8], [13] discussing the distribution of points in an n -dimensional Euclidean space E^n obtained from multiplicative congruential pseudo-random generators. For example, if x_0 , λ , β , μ , and n are given positive integers and if

$$(1) \quad x_i \equiv \lambda x_{i-1} \pmod{2^\beta} \quad (i = 1, 2, \dots)$$

or

$$(2) \quad x_i \equiv \lambda x_{i-1} + \mu \pmod{2^\beta} \quad (i = 1, 2, \dots),$$

Marsaglia [8] has shown that the points

$$(3) \quad (x_0, x_1, \dots, x_{n-1}), (x_n, \dots, x_{2n-1}), \dots$$

may lie on relatively few hyperplanes in E^n . This idea goes back at least to Franklin [4], [5]. Wood [14] describes a method he used to show that the points actually form a simple lattice in the case $n = 2$. Because his methods and results may be of some interest, it was thought that a report giving further details would be appropriate. In addition, a procedure for an extension to $3 \leq n \leq 6$ is given. Examples are presented for $2 \leq n \leq 6$. A more complete discussion of the general theory is given.

The interest in the present work arises from the need to know whether the generator produces points (3) which lie on few hyperplanes or lie on many. In the first instance, the pseudo-random points will not be uniformly distributed through the hypercube and hence the generator is probably not "good."

While this discussion bears some similarity to that of Coveyou and MacPherson [3], it has some advantages. First, it exhibits precisely the structure of the sets defined by (3). Secondly, it avoids a discussion of the Fourier analysis of lattice structure in which Coveyou and MacPherson couch their work. On the other hand, the present analysis has not been extended beyond $n = 6$ (but it is possible to extend the analysis), while Coveyou and MacPherson discuss $2 \leq n \leq 10$.

Received June 5, 1970.

AMS 1969 subject classifications. Primary 6515, 1061, 1063.

Key words and phrases. Multiplicative congruential pseudo-random numbers, Lehmer pseudo-random numbers, Monte Carlo integration, random numbers, lattices of random points, testing of random numbers, reduced cells, positive quadratic forms, quadratic forms.

* Work performed under the auspices of the U. S. Atomic Energy Commission.

Copyright © 1971, American Mathematical Society

In Section 8 the relation between the discrepancy theory of Zaremba (and others) and the present theory is discussed.

Applications of pseudo-random numbers in Monte Carlo calculations are well known. Applications in digital communications, especially in space communications, may not be so well known. See [6].

The computations were done on the Maniac II computer of the Los Alamos Scientific Laboratory. The Madcap language was used for the coding. This language can readily process arbitrarily large integers which is a requirement of our computations.

2. A Lattice in E^n . A lattice G_n in E^n consists of all vectors of the form $\mathbf{y} = \mathbf{e}_0 + \sum_{i=1}^n e_i \mathbf{y}_i$ where the \mathbf{e}_i ($1 \leq i \leq n$) are n fixed linearly independent vectors, the y_i are integers, positive, negative, or zero, and \mathbf{e}_0 is a fixed vector. (This definition is not standard in that the origin is not required to be in the lattice.) The $\{\mathbf{e}_i\}$ are said to form a basis of G_n . Put in other terms, a lattice G_n is a coset of a discrete subgroup H of the additive group of vectors in E^n where H has n linearly independent vectors. H is discrete if every $x \in H$ has a neighborhood free of points of H other than x . The basis vectors of G_n are then called generators of H .

Following van der Waerden [12, p. 276] one says the basis $\{\mathbf{e}_i\}$ is reduced (in the sense of Minkowski) if:

(1) \mathbf{e}_1 is the shortest (in the Euclidean norm) of all vectors $\sum_{i=1}^n e_i \mathbf{y}_i$ with the greatest common divisor of y_1, y_2, \dots, y_n : (y_1, \dots, y_n) , equal to 1,

(2) \mathbf{e}_k is the shortest of all vectors $\sum_{i=1}^n e_i \mathbf{y}_i$ with $(y_k, \dots, y_n) = 1$, for $k = 2, 3, \dots, n$.

Let N_1 be the length of the shortest nonzero vector $\mathbf{S}_1 = \sum_{i=1}^n e_i \mathbf{y}_i$. Let N_2 be the length of the shortest vector $\mathbf{S}_2 = \sum_{i=1}^n e_i \mathbf{y}_i$ which is linearly independent of \mathbf{S}_1 . And so on, one defines the successive minima N_1, N_2, \dots . Then if the $\{\mathbf{e}_i\}$ are reduced, it was shown by Mahler and Weyl [12] that

$$|\mathbf{e}_i| \leq \delta_i N_i, \quad i = 1, 2, \dots, n,$$

where $\delta_1 = 1$, $\delta_k = \max(1, \frac{1}{4}\delta_1 + \frac{1}{4}\delta_2 + \dots + \frac{1}{4}\delta_{k-1} + \frac{1}{4})$ for $k = 2, \dots, n$ and that

$$|\mathbf{e}_i| = N_i, \quad i = 1, 2, 3, 4,$$

where $|\mathbf{e}|$ denotes the Euclidean norm. It is this result which connects the reduced bases with the more intuitive idea of the "size" of the fundamental "cell" in the lattice and makes our theory a tool to study the distribution of pseudo-random points in the n -cube.

Minkowski [9] has stated the following for $n \leq 6$. $(\mathbf{e}_i)_{1 \leq i \leq n}$ is reduced if for every subset of $(\mathbf{e}_i)_{1 \leq i \leq n}$, say $(\mathbf{e}_i)_{1 \leq i \leq k}$, one has

$$|\mathbf{e}_{i_j}| \leq \left| \sum_{l=1}^k (\pm) C_l \mathbf{e}_{i_l} \right|, \quad j = 1, 2, \dots, k,$$

for all combinations of \pm signs and $(C_l)_{1 \leq l \leq k}$ ranging over the following values. If $k = 2, 3$, and 4 , $C_l = 1$. For $k = 5$, one of the C_l takes the values 1 and 2 and the remainder take the value 1. For $k = 6$, one of the C_l takes the values 1, 2, 3, another C_l takes the values 1 and 2, and the remainder take the value 1. (The cases $k = 5, 6$ are stated [9] without proof.)

The analysis given in van der Waerden [12] can be used to obtain algorithms for $n > 6$, but would not be optimal algorithms as is so for $2 \leq n \leq 6$.

For a set of vectors in E^n : $(a_i)_{1 \leq i \leq n}$, define $\det(a_i)$ to be the $n \times n$ determinant whose i th row consists of the components of the vector a_i . It is easy to show (see Cassels [2, p. 11, lines 7 to 13]) that for a given lattice G_n , a set $(a_i)_{0 \leq i \leq n}$ in G_n defines a basis $(a_i - a_0)_{1 \leq i \leq n}$ of G_n if and only if $0 < |\det(a_i - a_0)| \leq |\det(a'_i - a'_0)|$ for any other set $(a'_i)_{0 \leq i \leq n}$ in G_n such that $\det(a'_i - a'_0) \neq 0$. Further (see Cassels [2, pp. 9 and 10]), two sets $(a_i)_{0 \leq i \leq n}$ and $(a'_i)_{0 \leq i \leq n}$ in G_n both define a basis of G_n if and only if there exists an $n \times n$ matrix T with integer entries and with $\det T = \pm 1$ (unimodular matrix) so that $[a'_i - a'_0] = T[a_i - a_0]$, where $[a_i - a_0]$ is the $n \times n$ matrix whose i th row is the vector $a_i - a_0$.

A reduction algorithm is a procedure for obtaining from a basis of a lattice a reduced basis. Reduction algorithms are described and applied for $2 \leq n \leq 6$.

3. The Lattice Structure of Multiplicative Congruential Pseudo-Random Vectors. In this section n is an arbitrary positive integer. The following Lemmas 1 and 2 will be needed. They are taken from the book of Jansson [7, p. 68].

LEMMA 1. (1) *When $\lambda \equiv 3 \pmod{8}$ and $x_0 \equiv 1$ or $3 \pmod{8}$, each sequence produced by (1) is some permutation of all the numbers $8\nu + 1$ and $8\nu + 3$ ($\nu = 0, 1, \dots, 2^{\beta-3} - 1$).*

(2) *When $\lambda \equiv 3 \pmod{8}$ and $x_0 \equiv 5$ or $7 \pmod{8}$, each sequence produced by (1) is some permutation of all the numbers $8\nu + 5$ and $8\nu + 7$ ($\nu = 0, 1, \dots, 2^{\beta-3} - 1$).*

(3) *When $\lambda \equiv 5 \pmod{8}$ and $x_0 \equiv 1 \pmod{4}$, each sequence produced by (1) is some permutation of all the numbers $4\nu + 1$ ($\nu = 0, 1, \dots, 2^{\beta-2} - 1$).*

(4) *When $\lambda \equiv 5 \pmod{8}$ and $x_0 \equiv 3 \pmod{4}$, each sequence produced by (1) is some permutation of all the numbers $4\nu + 3$ ($\nu = 0, 1, 2, \dots, 2^{\beta-2} - 1$).*

Remark. A method of determining exactly what permutation occurs is illustrated by the following discussion.

Consider the case $\lambda \equiv 5 \pmod{8}$ and $x_0 = 1$. Denote the sequence generated by $\lambda = 5$ by

$$(4) \quad S_0 = \{x_i^{(0)}; i = 0, 1, 2, \dots, 2^{\beta-2} - 1\}$$

with $x_i^{(0)} \equiv 5^i \pmod{2^\beta}$. So every multiplier $\lambda \equiv 5 \pmod{8}$ with $0 < \lambda < 2^\beta$ occurs among the *odd* members of S_0 . Every $\lambda \equiv 5 \pmod{8}$, $0 < \lambda < 2^\beta$, multiplier has a representation of the form $\lambda = x_{2n+1}^{(0)} \equiv 5^{2n+1} \pmod{2^\beta}$. Let S_n be the sequence generated with $\lambda = x_{2n+1}^{(0)}$: $S_n = \{x_i^{(n)}; i = 0, 1, \dots, 2^{\beta-2} - 1\}$. One has

$$\begin{aligned} x_i^{(n)} &\equiv [x_{2n+1}^{(0)}]^i \pmod{2^\beta} [5^{2n+1}]^i \pmod{2^\beta} \\ &= 5^{(2n+1)i} \pmod{2^\beta} = x_{(2n+1)i}^{(0)}; \end{aligned}$$

i.e. when the multiplier $\lambda \equiv x_{2n+1}^{(0)}$ is used, the sequence obtained from $x_{i+1} \equiv \lambda x_i \pmod{2^\beta}$, $x_0 = 1$, consists of selecting every $(2n + 1)$ th number from (4), beginning with the first.

LEMMA 2. *If, in (2), $\lambda \equiv 1 \pmod{4}$ and $\mu \equiv 1 \pmod{2}$, then (2) produces a permutation of the numbers $0, 1, 2, \dots, 2^\beta - 1$.*

The following lemma is needed in the subsequent development.

LEMMA 3. *Let $A \subset E^n$ be a point set such that every point in A has integer co-*

ordinates and there are $n + 1$ vectors in A , say $e_i, 0 \leq i \leq n$, so that $e_i - e_0, 1 \leq i \leq n$, are linearly independent. Suppose for any $a_i \in A, 0 \leq i \leq n$, and any set of integers $k_i, 1 \leq i \leq n$, it is so that $a_0 + \sum_{i=1}^n k_i(a_i - a_0) \in A$. Then A is a lattice.

Proof. Choose $e_i, 0 \leq j \leq n$, in A so that $|\det(e_i - e_0)| = D$ has the least positive value where $1 \leq i \leq n$. By the hypothesis, every $e_0 + \sum_{i=1}^n k_i(e_i - e_0) \in A$ where the k_i are arbitrary integers. Let a be an arbitrary vector in A . Since $\{e_i - e_0; i = 1, 2, \dots, n\}$ is a linearly independent set, there exist real γ_i such that $a - e_0 = \sum_{i=1}^n \gamma_i(e_i - e_0)$. Suppose γ_j is not an integer for some $j, 1 \leq j \leq n$. Form

$$\begin{aligned} \left| \begin{array}{c} e_1 - e_0 \\ \vdots \\ e_{j-1} - e_0 \\ a - [\gamma_j](e_j - e_0) - e_0 \\ e_{j+1} - e_0 \\ \vdots \\ e_n - e_0 \end{array} \right| &= \left| \sum_{i=1}^n (\gamma_i - \delta_{ij}[\gamma_j])(e_i - e_0) \right| \\ &= \left| \begin{array}{c} e_1 - e_0 \\ \vdots \\ e_{j-1} - e_0 \\ (\gamma_j - [\gamma_j])(e_j - e_0) \\ e_{j+1} - e_0 \\ \vdots \\ e_n - e_0 \end{array} \right| = |\gamma_j - [\gamma_j]| D \end{aligned}$$

where $[\gamma_j]$ denotes largest integer in γ_j . Since $0 < |\gamma_j - [\gamma_j]| D < D$, it is false that D is the minimum positive value of $|\det(e_i - e_0)|$ where $e_i \in A$. Thus the γ_i are integers and therefore every $a \in A$ has a representation $a = e_0 + \sum_{i=1}^n \gamma_i(e_i - e_0)$ where the γ_i are integers. This completes the proof of the lemma. For a more general lemma, see Cassels [2, p. 78].

In Lemmas 4 to 8 the points defined by (1) or (2) and (3) or (1) or (2) and

$$(5) \quad (x_0, x_1, \dots, x_{n-1}), (x_1, \dots, x_n), (x_2, \dots, x_{n+1}), \dots$$

are discussed. We make the following convention: The point sets (3) and (5) are to be regarded as point sets G_n in E^n which are continued by periodicity throughout E^n ; i.e., if $(t_1, t_2, \dots, t_n) \in G_n$, then $(t_1 + h_1 2^\beta, t_2 + h_2 2^\beta, \dots, t_n + h_n 2^\beta) \in G_n$ for all positive, zero, and negative integers h_i .

Remark. It might be objected that points generated by (1) and (5) or (2) and (5) would make a poor random-point generator, since such points would be highly correlated over a short run. However, the points defined by (1) and (3) or (2) and (3)

are a reasonably sized subset for small n of those mentioned before and a discussion of the lattice structure of (5) gives information about the lattice structure of (3).

LEMMA 4. *If in (1) $\lambda \equiv 5 \pmod{8}$ and $x_0 \equiv 1 \pmod{4}$, then the point set G_n given by (5) forms a lattice in E^n .*

Proof. If $u_i, 1 \leq i \leq n$, are the unit coordinate vectors in E^n and $x \in G_n$, the vectors $x, x + 2^\beta u_i$ are $n + 1$ vectors in G_n for which $(x + 2^\beta u_i) - x$ are linearly independent. Let $x_{i_j}, 0 \leq j \leq n$, be $n + 1$ points of G_n , not necessarily distinct. Let $k_i, 1 \leq i \leq n$, be arbitrary integers. Recall that $x_{i_j} = (x_{i_j}, x_{i_{j+1}}, \dots, x_{i_{j+n}}) = (x_{i_j}, \lambda x_{i_j} + h_1 2^\beta, \lambda^2 x_{i_j} + h_2 2^\beta, \dots, \lambda^{n-1} x_{i_j} + h_{n-1} 2^\beta)$ for some integers h_i where x_{i_j} is the i_j th number generated by (1). Form

$$\begin{aligned} \mathbf{x} &= \mathbf{x}_{i_0} + \sum_{i=1}^n k_i (\mathbf{x}_{i_i} - \mathbf{x}_{i_0}) \\ &= \left\{ x_{i_0} + \sum_{i=1}^n k_i (x_{i_i} - x_{i_0}), \lambda \left[x_{i_0} + \sum_{i=1}^n k_i (x_{i_i} - x_{i_0}) \right] + h_1 2^\beta, \dots, \right. \\ &\qquad \qquad \qquad \left. \lambda^{n-1} \left[x_{i_0} + \sum_{i=1}^n k_i (x_{i_i} - x_{i_0}) \right] + h_{n-1} 2^\beta \right\} \end{aligned}$$

where h_i are again integers and the k_i are arbitrary integers. Since every x in the sequence generated by (1) under the hypothesis on λ and x_0 has the form $4\nu + 1, \nu = 0, 1, \dots, 2^{\beta-2} - 1$ (Lemma 1, Part 3), $x_{i_0} + \sum_{i=1}^n k_i (x_{i_i} - x_{i_0})$ has the form $4\nu + 1$ for some integer ν . But every number of this form can be expressed as $4\nu + 1 = 4\nu_1 + 1 + h2^\beta$ with $0 \leq \nu_1 \leq 2^{\beta-2} - 1$ and h, ν_1 as integers. Thus $\mathbf{x} \in G_n$. Lemma 3 can now be applied to give the conclusion of the theorem.

In a similar way, Lemmas 5 to 8 can be proved, using Lemmas 1, 2, and 3.

LEMMA 5. *If, in (1), $\lambda \equiv 5 \pmod{8}$ and $x_0 \equiv 3 \pmod{4}$, then the point set G_n given by (5) forms a lattice in E^n .*

LEMMA 6. *In (1), let $\lambda \equiv 3 \pmod{8}$ or $\lambda \equiv 5 \pmod{8}$ and let x_0 be odd. Then the set $(x_{2n}, x_{2n+1}), n = 0, 1, 2, \dots$, is a lattice in E^2 .*

LEMMA 7. *In (1), let $\lambda \equiv 3 \pmod{8}$. Let G_n be the set of points in (5) determined with $x_0 \equiv 1, \text{ or } 3 \pmod{8}$ and G'_n be the same set, but with $x_0 \equiv 5 \text{ or } 7 \pmod{8}$. Then $G_n \cup G'_n$ forms a lattice in E^n .*

LEMMA 8. *In (2), let $\lambda \equiv 1 \pmod{4}$ and $\mu \equiv 1 \pmod{2}$. Then the set of points in (5) form a lattice in E^n and the basis vectors of the lattice do not depend on μ . The sequence $(x_{2n}, x_{2n+1}), n = 0, 1, 2, \dots$, forms a lattice in E^2 .*

Remark 1. The points x_i (extended by our convention) defined by

$$x_i \equiv 3x_{i-1} \pmod{2^3}, \quad x_0 = 1,$$

do not form a lattice on the line.

Remark 2. The structure of sequences generated by other generators, such as

1. $x_{n+1} \equiv \lambda x_n + \mu \pmod{p^\beta}$ (p an odd prime),
2. $x_{n+1} \equiv \lambda x_n \pmod{10^\beta}$,
3. $x_{n+1} \equiv a_0 x_n + a_1 x_{n-1} + \dots + a_i x_{n-i} \pmod{p}$ (p a prime),
4. $x_{n+1} \equiv x_n + x_{n-1} \pmod{2^\beta}$,

is discussed in Jansson [7] and a theory analogous to that discussed here might be developable.

4. **Reduction Algorithm in the Case $n = 2$.** Let G_2 be a lattice with a basis (e_1, e_2) . Then, by the discussion in Section 2, if w is an integer, $(e_1, e_2 + we_1)$ is a basis of G_2 since

$$\begin{pmatrix} e_1 \\ e_2 + we_1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ w & 1 \end{pmatrix} \begin{pmatrix} e_1 \\ e_2 \end{pmatrix}$$

and the matrix $\begin{pmatrix} 1 & 0 \\ w & 1 \end{pmatrix}$ is unimodular. w is chosen to minimize $(e_2 + we_1)^2$. Hence w must satisfy $(e_2 + (w - 1)e_1)^2 \geq (e_2 + we_1)^2 \leq (e_2 + (w + 1)e_1)^2$ or

$$(6) \quad -\frac{e_1 \cdot e_2}{e_1^2} - \frac{1}{2} \leq w \leq -\frac{e_1 \cdot e_2}{e_1^2} + \frac{1}{2}.$$

In order to determine w uniquely, the right-hand inequality in (6) is replaced by $<$ to give

$$(7) \quad -\frac{e_1 \cdot e_2}{e_1^2} - \frac{1}{2} \leq w < -\frac{e_1 \cdot e_2}{e_1^2} + \frac{1}{2}.$$

Call the basis $(e_1, e_2 + we_1)$ thus determined (e_1, e'_2) . Replace the basis (e_1, e'_2) by a new basis $(e_1 + w'e'_2, e'_2)$ where w' is the unique integer determined by

$$(8) \quad -\frac{e_1 \cdot e'_2}{e_2'^2} - \frac{1}{2} \leq w' < -\frac{e_1 \cdot e'_2}{e_2'^2} + \frac{1}{2}.$$

The above procedure is then iterated until two successive minimizing integers of the form w and w' are zero. The resulting basis (\bar{e}_1, \bar{e}_2) is reduced since, from (7), $\bar{e}_1^2 \geq 2\bar{e}_1 \cdot \bar{e}_2 \geq -\bar{e}_1^2$ and, from (8), $\bar{e}_2^2 \geq 2\bar{e}_1 \cdot \bar{e}_2 \geq -\bar{e}_2^2$ and therefore

$$\bar{e}_1^2 \geq 2 |\bar{e}_1 \cdot \bar{e}_2| \quad \text{and} \quad \bar{e}_2^2 \geq 2 |\bar{e}_1 \cdot \bar{e}_2|$$

which implies that \bar{e}_1 and \bar{e}_2 are in length less than or equal to the length of the diagonals of the parallelograms which have \bar{e}_1, \bar{e}_2 as adjacent sides. The above algorithm must eventually terminate since for each pair of steps of the algorithm for which w and w' is not both zero, the vectors (e_1, e_2) with integer coordinates are replaced by a pair of vectors (e'_1, e'_2) with integer coordinates such that $|e'_1| \leq |e_1|$ and $|e'_2| \leq |e_2|$, with strict inequality in one of the cases.

5. **Reduction Algorithm in the Case $3 \leq n \leq 6$.** Assume $3 \leq n \leq 6$. Let $E = (e_1, e_2, \dots, e_n)$ be a set of basis vectors of a lattice G_n . Stage 1 of the reduction algorithm consists in successively replacing each pair of distinct vectors in E by a reduced pair, using the reduction algorithm for $n = 2$. This replacement defines a unimodular transformation from E to new set of vectors E' and hence E' is a basis of G_n . This operation is repeated until no further reduction by pairs is possible.

Stage 2 of the algorithm consists in examining for each k -tuple $(e_{i_j})_{1 \leq j \leq k}$ the vectors $\sum_{i=1}^k (\pm) C_i e_{i_j}$ where the values of C_i are described in Section 2. If it is found for some combination of \pm signs and C_i 's and for some e_{i_j} that

$$|e_{i_j}| > \left| \sum_{i=1}^k (\pm) C_i e_{i_j} \right|,$$

the vector e_{i_j} is replaced by the vector $e'_{i_j} = \sum_{i=1}^k (\pm) C_i e_{i_j}$ (the transformation from $E = (e_i)_{1 \leq i \leq n}$ to $E' = (e_1, e_2, \dots, e'_{i_j}, \dots, e_n)$ is unimodular). Stage 1 of the

algorithm is repeated on E' . The algorithm must terminate after a finite number of steps, since if a basis is altered by an operation in stages 1 or 2, the alteration consists in replacing a vector with integer coordinates with a shorter vector having integer coordinates.

Remark. In the initial consideration of the problem of finding reduced bases for $n > 2$, the search technique suggested by Coveyou and MacPherson [3] was considered with some modifications suggested in van der Waerden [12]. Without preliminary reduction it was found in a typical example (of the type discussed in this paper) about 10^{19} vectors would have had to be examined to find the shortest nonzero vector. Techniques used in crystallography were also considered (see Azároff and Buerger [1] and Roof [10]). Tests showed that these techniques were unsatisfactory for our purposes, due to lack of precision and the amount of search required.

6. Finding Bases for Multiplicative Congruential Pseudo-Random Points. To apply the reduction algorithm to the determination of reduced bases for pseudo-random points of the form (3) or (5), it is necessary to find a basis of these points. The method of finding a basis is illustrated by an example.

Consider the example of the generator (1) with $\lambda \equiv 5 \pmod{8}$ and $x_0 \equiv 1 \pmod{4}$. To find a set of basis vectors for G_n defined by (5), choose a set of $n + 1$ vectors in G_n as follows:

$$\begin{aligned} r_0 &= (1, \lambda, \lambda^2, \dots, \lambda^{n-1}), \\ r_i &= (4\alpha_i + 1, \lambda(4\alpha_i + 1) + h_i^1 2^\beta, \dots, \lambda^{n-1}(4\alpha_i + 1) + h_i^{n-1} 2^\beta), \\ & \qquad \qquad \qquad i = 1, 2, \dots, n, \end{aligned}$$

where $\alpha_i, h_i^k, i = 1, 2, \dots, n, k = 1, 2, \dots, n - 1$, are arbitrary integers. A calculation gives

$$\det (r_i - r_0) = 2^{2+(n-1)\beta} \begin{vmatrix} \alpha_1 & h_1^1 & \dots & h_1^{n-1} \\ \alpha_2 & \vdots & & \vdots \\ \vdots & \vdots & & \vdots \\ \alpha_n & h_n^1 & & h_n^{n-1} \end{vmatrix}$$

and $|\det (r_i - r_0)|$ has its minimum nonzero value when $\alpha_1 = h_2^1 = h_3^2 = \dots = h_n^{n-1} = 1$, the other determinant entries being zero. Thus a set of basis vectors of G_n is given by (with these choices for α_i and h_i^i)

$$\begin{aligned} r_1 - r_0 &= 4(1, \lambda, \lambda^2, \dots, \lambda^{n-1}), \\ r_i - r_0 &= (0, 0, \dots, 2^\beta, 0, \dots, 0), \quad i = 2, 3, \dots, n, \end{aligned}$$

where 2^β appears in the i th place.

7. Examples. Tables 1 to 5 present a few examples. It is hoped that the captions are self-explanatory, except for "figure of merit" defined for a reduced basis $\{y_i; i = 1, 2, \dots, n\}$ by

$$\text{figure of merit} = \frac{\min_{1 \leq i \leq n} |y_i|}{\max_{1 \leq i \leq n} |y_i|}.$$

TABLE 1

Reduced bases for lattices defined by (5) in n -dimensions for the traditional generator:
 $x_i \equiv 5^{15}x_{i-1} \pmod{2^{35}}$, $x_0 \equiv 1 \pmod{4}$.

$n = 2$

148728
 -439772
 177304
 399828

Figure of merit = 3.89=01
 Cosine of angle: 3.99=02

$n = 3$

1351376
 -1252052
 7001788
 3459216
 2632220
 31449484
 9668432
 -3330708
 -17029604

Figure of merit = 2.84=01
 Cosines of angles: -2.92=01 -1.23=01 1.08=01

$n = 4$

77710304
 -62468112
 90832932
 -8465924
 39043680
 75233072
 16799700
 13009868
 57634528
 -63711888
 -88091708
 25674588
 -27397792
 1670832
 -4290828
 -127412820

Figure of merit = 8.24=01
 Cosines of angles: -4.48=01 2.00=01 3.41=01 8.02=02 -2.23=02 -1.36=01

n = 5

25827680	237787616	107896928	52851936	-186499232
192251352	157110264	-62209128	69997240	225757016
200526964	-247943196	45409428	18094468	-221466188
30054476	-48263716	-197075412	277932604	132652812
-40032076	34520868	-449544420	-183833916	-5688844

Figure of merit = 6.66-01
 Cosines of angles: -2.46-02 -5.19-02 -3.48-01 -3.14-01 -3.81-01 4.64-01
 6.78-02 -1.06-01 -2.04-01 1.87-01

n = 6

-305139696	650889424	269544080	478429520	394774800	49237456
-106314720	-15850080	-224044768	538719904	-146950624	59535264
428761576	-6351672	20150824	548909064	776354920	-961483448
140595992	-318464968	-278002984	-35706632	598755480	453202872
-625061744	-31756464	-370532592	341922864	-235736688	-452912048
-173820160	-715470080	98881280	286065408	285583104	129760000

Figure of merit = 4.33-01
 Cosines of angles: 2.71-01 2.80-01 -9.67-02 -2.22-01 -1.54-01 8.54-02
 -5.19-02 -4.54-02 2.43-01 5.34-02 -1.56-01 1.53-01
 -3.71-01 5.29-01 -1.59-01

TABLE 2
 Reduced bases for lattices defined by (5) in n -dimensions for a "bad" generator:
 $x_i \equiv 5x_{i-1} \pmod{2^{25}}, x_0 \equiv 1 \pmod{4}$.

	n = 2		
	h	20	
	-6607641992	1321528408	
Figure of merit =	3.03-09		
Cosine of angle =	1.40-09		
	n = 3		
	h	100	
	-6861391488	52779928	
	-1319498400	-6597492000	
Figure of merit =	1.49-08		
Cosines of angles =	-5.98-09	4.57-09 1.92-01	
	n = 4		
	h	20	500
	-6871525460	2111068	52776700
	-1372194024	-6860970120	10555340
	-263883468	-1319417340	54887768
		-6597086700	1374304868
Figure of merit =	7.29-08		
Cosines of angles =	1.00-09	5.12-09 -3.33-08	1.96-01 3.84-02 1.96-01

$n = 5$
 u 20 100 500 2500
 -6871930788 84428 422140 2110700 10553500
 -1374301716 6871508580 2195468 10977340 54886700
 274438132 1372190660 6860953300 54971868 -274859340
 -52776564 -263882820 -1319414100 6597070500 1374385868
 Figure of merit = 3.64=07
 Cosines of angles: -2.69=07 -1.35=07 1.46=07 -4.29=08 1.96=01 -3.92=02
 7.68=03 -2.00=01 3.92=02 -1.96=01

$n = 6$
 u 20 100 500 2500
 6871946996 -3388 -16940 84700 423500
 1374386020 6871930100 -87868 439340 2196700
 -274860316 -1374301580 -6871507900 2198868 10994340
 54887620 274438100 1372190500 6860952500 54975868
 -10555312 -52776560 -263882800 -1319414000 6597070000
 Figure of merit = 1.82=06
 Cosines of angles: -9.56=07 -8.70=07 -5.29=08 -3.11=07 -1.53=07 1.96=01
 -3.92=02 7.83=03 -1.54=03 -2.00=01 3.99=02 -7.83=03
 -2.00=01 3.92=02 -1.96=01

TABLE 3

Reduced bases for lattices defined by (5) in n -dimensions for the "shift and add" generator:
 $x_i \equiv (2^{17} + 5)x_{i-1} \pmod{2^{35}}$, $x_0 \equiv 1 \pmod{4}$.

$n = 2$

262136
 -262132

Figure of merit = 1.00 00
 Cosine of angle: 7.63-05

$n = 3$

262140
 -786416
 -1183269268

786412
 -1834928
 473151516

1310620
 1311120
 -46491508

Figure of merit = 1.22-03
 Cosines of angles: 1.67-02 4.91-04 4.55-04

$n = 4$

-262140
 1572844
 486529588
 -1356543000

-786412
 5242780
 1164395268
 59243400

-1310620
 13106700
 -519287020
 146270632

6554100
 -2500
 56986468
 -18378680

Figure of merit = 4.93-03
 Cosines of angles: -2.27-01 1.18-03 -3.42-04 3.28-03 4.85-03 -3.58-01

n = 5

786424	2883544	9174840	19659800	-32773000
544986936	1356543000	-59243400	-146270632	18378680
148109424	486791728	1165181680	-517976400	50432368
-1310708	-4980676	-17039060	-45873700	-32760500
1373057856	-4980672	-16514752	-40630720	6561600

Figure of merit = 2.68-02

Cosines of angles: 5.50-03 6.91-03 2.25-05 -1.85-03 3.71-01 -4.07-03
 3.71-01 -4.47-03 1.08-01 1.91-03

n = 6

524284	2097132	7864220	26213900	65533500	-12500
-786424	-2883544	-9174840	-19659800	32773000	819225000
1373844280	-2097128	-7339912	-20970920	-26211400	262159000
-549443340	-1373057852	5504980	21757732	79952820	255584900
-163313628	-543414092	-1351300220	72350100	146268132	-346071180
39845500	148371564	487578140	1166492300	-524530500	-47874132

Figure of merit = 4.69-02

Cosines of angles: 2.65-02 -1.62-02 2.51-02 -3.40-03 3.79-03 1.86-01
 1.75-01 -2.14-01 -7.46-02 -3.27-01 -1.46-01 1.43-02
 3.31-01 -1.22-01 -3.46-01

TABLE 4
 Reduced bases for lattices defined by (5) in n -dimensions for a generator with a "randomly selected" multiplier:
 $x_i \equiv \lambda x_{i-1} \pmod{2^{29}}$ where $\lambda^n = 273673163157 \equiv 5 \pmod{8}$ and $x_0 \equiv 1 \pmod{4}$.

	$n = 2$		
		-83308	
		394836	
		-365532	
		82660	
Figure of merit	$= 9.29=01$		
Cosine of angle:	$= 4, 17=01$		
	$n = 3$		
		19835632	24341936
		-3506608	9360528
		-5256076	1684356
			19367792
			-19970096
			2441172
Figure of merit	$= 1.64=01$		
Cosines of angles:	$= 2, 77=01$	$= 7, 18=02$	$= 1, 08=01$
	$n = 4$		
		-49630292	18153244
		62445036	19622620
		3435528	9229896
		-46485684	2152592
			-91924660
			6559628
			-75321208
			104324060
Figure of merit	$= 5.13=01$		
Cosines of angles:	$= 3, 15=01$	$= 2, 75=01$	$= 4, 55=01$
		$= 2, 97=01$	$= 1, 17=01$
			$= 5, 92=02$

n = 5

246340616	99260584	183849320	1591688
-83006312	203312248	-50576712	75042072
-155976676	92658660	106194884	109603092
-258676968	29735252	179924404	-307638844
-150242428	267746004	-149567436	-202573244

Figure of merit = 6.86-01

Cosines of angles: 1.90-02 -1.65-01 -2.06-01 -2.93-01 -2.12-01 -9.47-02
 1.45-01 2.44-01 4.15-01 4.49-01

n = 6

-159068752	-497660176	325706160	122489712	235542576	-122993680
233248740	-491058252	1622980	20044448	127531172	495229812
533124172	-178348996	535620332	-408262820	262620300	-155963012
76404388	120770932	129014916	158588116	-693265564	411112756
-164957284	-502716980	115444668	-628502420	-575046948	37297676
-491058252	1622980	20044448	127531172	495229812	498413180

Figure of merit = 6.77-01

Cosines of angles: 3.77-01 3.21-01 -3.91-01 1.43-01 3.59-01 1.22-01
 1.62-01 3.58-02 3.21-01 -2.84-01 1.74-01 -1.87-01
 3.00-01 -1.73-01 -2.71-01

TABLE 5

Reduced bases for lattices defined by (5) in n -dimensions for the generator: $x_i \equiv (2^{13} + 5)x_{i-1} \pmod{2^{47}}$.
 Note the unusual structure for $n = 5$, two basis vectors being $(0, 2^{47}, 0, 0, 0)$ and $(0, 0, 2^{47}, 0, 0)$.

$n = 2$

4 32788
 -17169389564 2099220

Figure of merit = 1.91-06
 Cosine of angle: 2.69-07

$n = 3$

4 32788 268763236
 -17169389820 788 6459236
 -2094596 -17169403412 -111412836

Figure of merit = 1.57-02
 Cosines of angles: 3.76-04 6.61-03 1.20-04

$n = 4$

-17169389820 788 6459236 52946357492
 34342969340 34342969324 34342838172 33267785228
 -6283272 -51503980584 34336218936 -20990092264
 -34334590196 34340874044 17167828012 -12802140964

Figure of merit = 7.79-01
 Cosines of angles: 3.09-01 -3.05-01 -2.99-02 -2.89-01 4.53-02 -2.62-01

n = 5

```

40      327880      2687632360 22030522454920 17995003093416
-0 140737488355328
 0      140737488355328
220      1803340      14781977980 -19569614853268 28603772836124
-12     -98364      -806289708 -6609156736476 8675247907508
    
```

Figure of merit = 7.75=02

Cosines of angles: 1.15=06
 -9.02=09 4.27=04 7.39=05 9.99=01
 9.45=05 8.48=02 3.39=02 0.0 5.20=08

n = 6

```

360576037316 154522417364 -17140065244 240373550260 17154506628 121997525612
-154398829144 1030192704584 240298155368 -600857423096 583795318568 295622220744
-669578975812 223181126828 -171652010148 343356370124 -257601200132 -494712152084
326224686276 -51474653140 274733277404 188861195340 -17153706628 128555125612
-188802537120 497975135968 515027225184 -446485826592 -669623336096 -140440121120
 85844862172 -17106552756 515075414396 -51476855828 240283843868 -718168788596
    
```

Figure of merit = 3.43=01

Cosines of angles: 1.07=01
 -3.27=02 3.27=01 4.04=02 2.25=01 2.20=01 1.40=01
 1.61=01 1.63=02 2.00=01
 -1.41=01 -4.35=01 -1.93=01 2.25=01 -1.55=01
 3.47=01 2.20=01 1.40=01

The figure of merit provides some measure (neglecting angle) of the departure of the reduced cell from "squareness." $x.xx - 0a$ means $x.xx \cdot 10^{-a}$. The angles refer to angles between the edges. Angles between higher-dimensional flats have not been calculated, but it might be useful to do so. Each row of the tables lists the components of a vector.

It is seen from these tables that multipliers of simple structure, such as 5 or $2^{17} + 5$ produce lattices of points which depart greatly from a uniform distribution throughout the cube. Multipliers of more complex structure, such as 5^{15} or the "randomly" selected multiplier 273673163157 produce better lattices. A typical time for the computation of a table on Maniac II was 20 minutes. The CDC 6600 is perhaps 8 times faster than Maniac II.

It is hoped that more examples with a more complete discussion can be presented in the future.

8. Connection Between Lattices of Pseudo-Random Points and the Theory of Discrepancy. Zaremba [15], Schmidt [11], and others (see references in [11] and [15]) have discussed the notion of the discrepancy $D(S)$ of a finite number of points S in the unit cube $I_n \subset E^n$. The quantity $D(S)$ can be used to estimate the error in the evaluation of a multidimensional integral. As an example, if $f(x, y)$ is of bounded variation in the sense of Hardy and Krause over I_2 and $S = \langle x_0, \dots, x_{N-1} \rangle$ is an arbitrary sequence of points in I_2 , then

$$\left| \int_{I_2} f(\mathbf{x}) \, d\mathbf{x} - N^{-1} \sum_{k=0}^{N-1} f(\mathbf{x}_k) \right| \leq V^2(f)D(S) + V(f(x, 1))D(X) + V(f(1, y))D(Y),$$

where V and V^2 denote one- and two-dimensional variation and X and Y are projections of S on the x and y axis respectively. Roth (see [15]) has proved that, for E^n ,

$$D(S) \geq C_n N^{-1} (\log N)^{(n-1)/2}$$

for some constant C_n .

It seems to be a reasonable conjecture that if the figure of merit (see Section 7) of a lattice G_n is very small, then $D(G_n \cap I_n)$ is large. Conversely, if the figure of merit is near 1, $D(G_n \cap I_n)$ is small.

It should be remarked that the application of this lattice theory to much shorter segments of the full period of the generator sequence depends on the extent to which the lattice properties are reflected in the segments.

Acknowledgment. The authors thank Dr. W. W. Wood of this laboratory for valuable help and suggestions in the preparation of this paper.

Postscript. After completion of the above paper, the following important paper came to the authors' attention:

R. R. Coveyou, "Random number generation is too important to be left to chance," *Studies in Appl. Math.*, v. 3, 1970, pp. 70-111.

That paper has things in common with our paper. However, our paper was written independently and differs from the former in important details. It was thought best to not revise the present paper.

University of California
Los Alamos Scientific Laboratory
Los Alamos, New Mexico 87544

1. L. V. AZÁROFF & M. J. BUERGER, *The Powder Method in X-Ray Crystallography*, McGraw-Hill, New York, 1958.
2. J. W. S. CASSELS, *An Introduction to the Geometry of Numbers*, Die Grundlehren der math. Wissenschaften in Einzeldarstellungen mit besonderer Berücksichtigung der Anwendungsgebiete, Band 99, Springer-Verlag, Berlin, 1959. MR 28 #1175.
3. R. R. COVEYOU & R. D. MACPHERSON, "Fourier analysis of uniform random number generators," *J. Assoc. Comput. Mach.*, v. 14, 1967, pp. 100–119. MR 36 #4779.
4. J. N. FRANKLIN, "Deterministic simulation of random processes," *Math. Comp.*, v. 17, 1963, pp. 28–59. MR 26 #7125.
5. J. N. FRANKLIN, "Equidistribution of matrix-power residues modulo 1," *Math. Comp.*, v. 18, 1964, pp. 560–568.
6. S. W. GOLOMB, L. D. BAUMERT, M. F. EASTERLING, J. J. STIFFLER & A. J. VITERBI, *Digital Communications with Space Applications*, Prentice-Hall, Englewood Cliffs, N. J., 1964.
7. B. JANSSON, *Random Number Generators*, Almqvist & Wiksell, Stockholm, 1966. MR 36 #7297.
8. G. MARSAGLIA, "Random numbers fall mainly in the planes," *Proc. Nat. Acad. Sci. U.S.A.*, v. 61, 1968, pp. 25–28. MR 38 #3998.
9. H. MINKOWSKI, "Zur Theorie der positiven quadratischen Formen," *J. Reine Angew. Math.*, v. 101, 1887, pp. 196–202.
10. R. B. ROOF, JR., *A Theoretical Extension of the Reduced-Cell Concept in Crystallography*, Los Alamos Scientific Laboratory Report, LA-4038, 1969.
11. W. SCHMIDT, "Irregularities of distribution. IV," *Invent. Math.*, v. 7, 1969, pp. 55–82. MR 39 #6838.
12. B. L. VAN DER WAERDEN, "Die Reduktionstheorie der positiven quadratischen Formen," *Acta Math.*, v. 96, 1956, pp. 265–309. MR 18, 562.
13. P. H. VERDIER, "Relations within sequences of congruential pseudo-random numbers," *J. Res. Nat. Bur. Standards Sect. B*, v. 73B, 1969, pp. 41–44. MR 39 #1081.
14. W. W. WOOD, "Monte Carlo calculations for hard disks in the isothermal-isobaric ensemble," *J. Chem. Phys.*, v. 48, 1968, pp. 415–434.
15. S. K. ZAREMBA, "The mathematical basis of Monte Carlo and quasi-Monte Carlo methods," *SIAM Rev.*, v. 10, 1968, pp. 303–314. MR 38 #1810.