

# A Continued Fraction Algorithm for Real Algebraic Numbers\*

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**Abstract.** Let  $\alpha$  denote a real algebraic number that is a root of a polynomial  $f(x) \in \mathbf{Z}[x]$ . The purpose of this paper is to state an algorithm for finding the simple continued fraction expansion of  $\alpha$ . Furthermore, an application of the algorithm to sign determination in real algebraic number fields is given.

**1. Introduction.** The task of constructively computing the simple continued fraction expansion (see [2]) for a real root  $\alpha$  of a polynomial

$$f(x) = b_0x^m + b_1x^{m-1} + \cdots + b_m \quad (b_0 \neq 0)$$

over the rational integers  $\mathbf{Z}$  raises no essential difficulties provided that  $\alpha$  is the sole real root of  $f(x)$ . However, if  $f(x)$  happens to have more than one real root, the problem arises of discriminating between the continued fraction expansion of  $\alpha$  and of the other real roots of  $f(x)$ .

An attempt to solve this problem was made by Zassenhaus [5], who showed that, after a finite number of steps, the so-called "reduced state" of the continued fraction expansion of  $\alpha$  (see below) is reached. (See also [2].) From there on, the discrimination of the real roots is automatically guaranteed. Unfortunately, no indication is given in Zassenhaus' method as to when the reduced state will be attained for a given  $\alpha$ , nor does there seem to exist a simple device for achieving that state (cf. [6]). Nonetheless, a computer program was written by Smith [3] in which the method is applied to some special cases.

**2. The Continued Fraction Algorithm.** In this paper, we describe a different continued fraction algorithm that furnishes a solution to the discrimination problem mentioned above and that, moreover, appears to be much simpler than the routine designed by Zassenhaus [5].

Let us first remark that, as Zassenhaus [5] observed, it is expedient to reduce  $f(x)$  to a polynomial having no multiple factors. We can eliminate them by replacing  $f(x)$  by the polynomial  $f(x)/(f(x), f'(x))$ . In the trivial case in which  $\alpha$  is a *rational* root of  $f(x)$ , the algorithm will simply terminate after a finite number of steps.

We confine ourselves therefore to giving a description of the algorithm as applied to an *irrational* real root  $\alpha$  of the (not necessarily irreducible) polynomial  $f(x)$  in  $\mathbf{Z}[x]$ .

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The polynomial  $f(x)$  may, moreover, be supposed to have no rational roots at all. The continued fraction expansion of  $\alpha$  is then calculated assuming that  $\alpha$  is isolated by rational numbers (or infinity)  $r$  and  $s$ ; i.e.,  $\alpha$  is the unique root of  $f(x)$  in the closed interval  $[r, s]$ . Put  $r_0 = r, s_0 = s$ , and define the 0th successor  $\alpha_0$  of  $\alpha$  by

$$\alpha_0 = \alpha,$$

the 0th partial denominator  $a_0$  of  $\alpha$  by

$$a_0 = [\alpha_0],$$

where  $[ ]$  designates the greatest integer function, and the 0th successor polynomial  $f_0(x)$  of  $f(x)$  by

$$f_0(x) = f(x).$$

We have  $f_0(\alpha_0) = 0$ .

Let us assume by induction that, for an integer  $n \geq 1, \alpha_{n-1}$  is an irrational real root of a polynomial

$$f_{n-1}(x) = b_{0,n-1}x^m + b_{1,n-1}x^{m-1} + \dots + b_{m,n-1} \quad (b_{0,n-1} \neq 0), (b_{i,0} = b_i \text{ for } 0 \leq i \leq m),$$

over  $\mathbf{Z}$  having neither multiple factors nor rational roots, and that  $\alpha_{n-1}$  is the unique root of  $f_{n-1}(x)$  in the closed interval  $[r_{n-1}, s_{n-1}]$ .

Next, put  $a_{n-1} = [\alpha_{n-1}]$ , and let

$$\begin{aligned} r_n &= (s_{n-1} - a_{n-1})^{-1} && \text{if } s_{n-1} < a_{n-1} + 1, \\ &= 1 && \text{otherwise,} \\ s_n &= (r_{n-1} - a_{n-1})^{-1} && \text{if } r_{n-1} > a_{n-1}, \\ &= \infty && \text{otherwise.} \end{aligned}$$

Define the  $n$ th successor  $\alpha_n$  of  $\alpha$  by

$$\alpha_n = (\alpha_{n-1} - a_{n-1})^{-1},$$

the  $n$ th partial denominator  $a_n$  of  $\alpha$  by

$$a_n = [\alpha_n],$$

and the  $n$ th successor polynomial  $f_n(x)$  of  $f(x)$  by

$$f_n(x) = x^m f_{n-1}(x^{-1} + a_{n-1}).$$

Clearly,  $f_n(x)$  is a polynomial over  $\mathbf{Z}$  having neither multiple factors nor rational roots, and  $\alpha_n$  is one of the irrational real roots of  $f_n(x)$ . Moreover,  $\alpha_n$  is the unique root of  $f_n(x)$  in the closed interval  $[r_n, s_n]$ . Note that for  $n \geq 1$ , we have  $\alpha_n > 1$  and  $1 \leq r_n < s_n \leq \infty$ .

The definition of  $r_n, s_n$  and  $\alpha_n$  leads us to the following observation which is of significance for the discrimination problem mentioned at the beginning (cf. [2] and the Theorem of Vincent [4]).

**THEOREM.** *Under the above hypothesis on  $\alpha$ ,  $r$ , and  $s$ , there exists  $n_1$  such that, for all  $n \geq n_1$ , we have*

$$r_n = 1 \quad \text{and} \quad s_n = \infty.$$

*Proof.* The assertion results from two facts that are immediate consequences of the definition of  $r_n$ ,  $s_n$  and  $\alpha_n$ .

(1) If  $r_n = 1$  or  $s_n = \infty$  for some integer  $n \geq 1$ , then it follows that

$$r_{n+i} = 1 \text{ for all even or all odd natural numbers } i, \\ \text{respectively,}$$

and that

$$s_{n+j} = \infty \text{ for all odd or all even natural numbers } j, \\ \text{respectively.}$$

(2) For all integers  $n \geq 1$ , the following hold:

$$\text{either } r_n = 1 \text{ or } a_{n-1} = [s_{n-1}]$$

and

$$\text{either } s_n = \infty \text{ or } a_{n-1} = [r_{n-1}].$$

Once we have arrived at an index  $n_1 \geq 1$  such that  $r_{n_1} = 1$  and  $s_{n_1} = \infty$ , statement (1) implies that  $r_n = 1$  and  $s_n = \infty$  for all  $n \geq n_1$ . To see that  $n_1$  exists, consider the sequences  $S = \{[r_0], [s_1], [r_2], \dots\}$  and  $T = \{[s_0], [r_1], [s_2], \dots\}$ , which are initially the continued fraction expansions of  $r$  and  $s$ , respectively. By (2),  $S$  and  $T$  must each eventually differ from the continued fraction expansion  $\{a_0, a_1, a_2, \dots\}$  of  $\alpha$  since  $r \neq \alpha$  and  $s \neq \alpha$ .  $S$  and  $T$  each then become  $\{\dots, 1, \infty, 1, \infty, \dots\}$ .

The actual determination of the partial denominators  $a_n$  of  $\alpha$  can now be carried through in the following manner (see also [5]).

First, we find improved bounds for the irrational real root  $\alpha_n$  of  $f_n(x)$  where  $n \geq 0$ . To this end, we have to introduce the set

$$\alpha_n = \alpha_n^{(1)}, \alpha_n^{(2)}, \dots, \alpha_n^{(m)}$$

of the complex roots of  $f_n(x)$ . We recall that these roots can be defined inductively by setting  $\alpha_0^{(i)} = \alpha^{(i)}$  and, for  $n \geq 1$ ,

$$\alpha_n^{(i)} = (\alpha_{n-1}^{(i)} - a_{n-1})^{-1} \quad (i = 1, 2, \dots, m).$$

Also, we use the  $n$ th convergent of the continued fraction expansion of  $\alpha$ , that is, the fraction (see [2])

$$[a_0, a_1, \dots, a_n] = p_n/q_n \quad (p_n, q_n \in \mathbf{Z}).$$

As usual, define  $p_{-1} = 1$  and  $q_{-1} = 0$ .

The integers  $p_{n-1}$ ,  $p_{n-2}$  and  $q_{n-1}$ ,  $q_{n-2}$  ( $n \geq 1$ ) appear in the formula connecting  $\alpha^{(i)}$  with  $\alpha_n^{(i)}$ , namely,

$$\alpha^{(i)} = (p_{n-1}\alpha_n^{(i)} + p_{n-2}) / (q_{n-1}\alpha_n^{(i)} + q_{n-2}) \quad (i = 1, 2, \dots, m)$$

or, conversely,

$$\alpha_n^{(i)} = -(q_{n-2}\alpha^{(i)} - p_{n-2}) / (q_{n-1}\alpha^{(i)} - p_{n-1}).$$

We write the latter relation for  $n \geq 2$  in the form

$$\alpha_n^{(i)} = \frac{\alpha^{(i)} - p_{n-2}/q_{n-2} \cdot q_{n-2}}{\alpha^{(i)} - p_{n-1}/q_{n-1} \cdot q_{n-1}} \quad (i = 1, 2, \dots, m).$$

Noting that  $p_{n-2}/q_{n-2} \rightarrow \alpha$  and  $p_{n-1}/q_{n-1} \rightarrow \alpha$ , as  $n \rightarrow \infty$ , and that  $\alpha \neq \alpha^{(i)}$  for all  $i$  in the interval  $1 < i \leq m$ , we conclude that, for  $i \neq 1$  and for all large  $n$ , the  $|\alpha_n^{(i)}|$  are asymptotic to  $q_{n-2}/q_{n-1}$ . On the other hand, it follows from the second of the two relations

$$\begin{aligned} p_{n-1} &= p_{n-2}a_{n-1} + p_{n-3} \\ q_{n-1} &= q_{n-2}a_{n-1} + q_{n-3} \end{aligned} \quad (n \geq 2),$$

or, respectively, from the definition of  $q_{-1}$  and  $q_0$  that

$$q_{n-2}/q_{n-1} \leq a_{n-1}^{-1} \quad (n \geq 2)$$

with strict inequality for  $n \geq 3$ . For all large  $n$ , the conjugates  $\alpha_n^{(i)}$  of  $\alpha_n = \alpha_n^{(1)}$  satisfy

$$|\alpha_n^{(i)}| < a_{n-1}^{-1} \quad (1 < i \leq m).$$

It is clear from the above relations that there exists  $n_2$  such that, for all  $n \geq n_2$ , the following two conditions are fulfilled:

$$\begin{aligned} \alpha_n &> 1, \\ 0 &< -\operatorname{Re}(\alpha_n^{(i)}) \leq |\alpha_n^{(i)}| < 1 \quad (1 < i \leq m), \end{aligned}$$

where “Re” denotes the real part of a complex number. This is what Zassenhaus [5] calls the *reduced state* of the continued fraction expansion of  $\alpha$ . Thus, for all large  $n$ ,  $\alpha_n$  is a *PV* number.

As soon as the reduced state is reached, we know, because of the relation

$$\sum_{i=1}^m \alpha_n^{(i)} = -b_{1,n}/b_{0,n}$$

on the roots  $\alpha_n^{(i)}$  of  $f_n(x)$ , that  $\alpha_n = \alpha_n^{(1)}$  lies in the interval

$$-b_{1,n}/b_{0,n} < \alpha_n < (m - 1) - b_{1,n}/b_{0,n}.$$

The upper bound for  $\alpha_n$  can be further improved. Specifically, from the relations derived above, we infer that  $\alpha_n$  is asymptotic to  $(m - 1)q_{n-2}/q_{n-1} - b_{1,n}/b_{0,n}$  and moreover, that there exists  $n_3$  such that, for all  $n \geq n_3$ , we have

$$\alpha_n < (m - 1)/a_{n-1} - b_{1,n}/b_{0,n}.$$

Now, if  $n \geq 1$  and  $a_0, a_1, \dots, a_{n-1}$  are already computed, we calculate  $a_n$  via a modified binary search process in the interval  $u_n \leq a_n \leq v_n$  which is roughly defined as follows. Put

$$n_4 = \max\{n_1, n_2, n_3\},$$

where  $n_i$  are the preceding index bounds. Then, we put for  $n < n_4$ ,

$$\begin{aligned} u_n &= [r_n] \quad \text{if } n \text{ is even,} \\ &= [s_n], \quad \text{if } n \text{ is odd,} \end{aligned}$$

$$v_n = \min\{[s_n], [t_n]\}, \text{ if } n \text{ is even,}$$

$$= \min\{[r_n], [t_n]\}, \text{ if } n \text{ is odd,}$$

where

$$t_n = 1 + \max_{1 \leq i \leq m} \{ |b_{i,n}| / |b_{0,n}| \},$$

and, for  $n \geq n_4$ ,

$$u_n = \max\{1, [-b_{1,n}/b_{0,n}]\},$$

$$v_n = [(m - 1)/a_{n-1} - b_{1,n}/b_{0,n}].$$

Note that, for  $n \geq 1$ ,  $u_n$  and  $v_n$  are positive integers.

The  $n$ th partial denominator  $a_n$  of  $\alpha$  is then determined as the unique natural number  $\lambda_n$  in the interval  $u_n \leq \lambda_n \leq v_n$  for which

$$\text{sgn } f_n(\lambda_n) \neq \text{sgn } f_n(\lambda_n + 1).$$

Before describing the binary search process for  $a_n$ , we note that, if  $n \geq n_4$ , it is expedient to precede the binary search with the sign test for

$$\lambda_n = [(m - 1)q_{n-2}/q_{n-1} - b_{1,n}/b_{0,n}],$$

because the number in square brackets is, as we have seen, a good approximation to  $\alpha_n$ . This, of course, requires computation of the  $q_n$ . If  $\text{sgn } f_n(\lambda_n) \neq \text{sgn } f_n(\lambda_n + 1)$  for this  $\lambda_n$ , then  $a_n = \lambda_n$ . Otherwise, we start the binary search as follows. We put  $\lambda_n = v_n$  and check whether  $\text{sgn } f_n(\lambda_n) \neq \text{sgn } f_n(\lambda_n + 1)$ . If so, then  $a_n = v_n$ . If not, we know that  $u_n \leq a_n \leq v_n - 1$ . Unless  $u_n = v_n - 1$ , in which case  $a_n = u_n$ , we put

$$w_n = \lfloor \frac{1}{2}(u_n + v_n) \rfloor$$

and compare the signs of  $f_n(w_n)$  and  $f_n(v_n)$ , say. If they differ, we replace  $u_n$  by  $w_n$ ; otherwise, we leave  $u_n$  unchanged and substitute  $w_n$  for  $v_n$ . The search process is then repeated (if need be) with respect to the new interval, until  $u_n = v_n - 1$ .

This algorithm has been implemented as a computer program which we shall use to build the example of Section 4.

**3. An Application of the Algorithm to Sign Determination.** In this section, we shall outline a method for performing sign determination in a real algebraic number field

$$K = \mathbb{Q}(\alpha)$$

over the field of rational numbers  $\mathbb{Q}$ , where  $\alpha$  is an irrational real root of a (not necessarily irreducible) polynomial  $f(x)$  in  $\mathbb{Z}[x]$  of degree  $m > 1$  as before. This method seems to be somewhat simpler than the one proposed by Kempfert [1] and Zassenhaus [6]; however, their method applies to any ordered field.

Every element  $\beta \in K$  can be represented in the form  $\beta = g(\alpha)$  with a polynomial  $g(x)$  in  $\mathbb{Q}[x]$  of degree  $< m$ .

First of all, we may assume that  $g(\alpha) \neq 0$ , since if  $g(\alpha)$  were 0, then  $(f(x), g(x)) \neq 1$ .

To determine the sign of  $g(\alpha)$ , we employ the continued fraction algorithm of Section 2 in order to approximate  $\alpha$  by its convergents  $p_n/q_n$ . The theory of continued

fractions yields, for the approximation of  $\alpha$  by  $p_n/q_n$ , the estimate (see [2])

$$|\alpha - p_n/q_n| < 1/q_n^2,$$

where  $q_n \rightarrow \infty$ , as  $n \rightarrow \infty$ .

We shall show that, for all large  $n$ , the sign of  $g(\alpha)$  can be obtained from the relation

$$\text{sgn } g(\alpha) = \text{sgn } g(p_n/q_n).$$

To this end, we note that, by the mean value theorem (cf. [6]), the formula

$$g(\alpha) - g(p_n/q_n) = g'(\xi)(\alpha - p_n/q_n)$$

is valid, where  $g'(x)$  denotes the derivative of  $g(x)$  and  $\xi$  is a real number lying between  $\alpha$  and  $p_n/q_n$ . Let  $M$  be a bound for  $g'(x)$  for  $x$ , say between  $p_0/q_0$  and  $p_1/q_1$ . We thus have

$$|g(\alpha) - g(p_n/q_n)| < M/q_n^2.$$

Then  $g(p_n/q_n) \rightarrow g(\alpha)$ . For large enough  $n$ ,  $|g(p_n/q_n)| \geq M/q_n^2$  and then  $\text{sgn } g(p_n/q_n) = \text{sgn } g(\alpha)$ .

**4. An Example for the Continued Fraction Algorithm.** We compute here the continued fraction expansion for three roots of the polynomial

$$f(x) = x^7 - 7x + 3,$$

which has three irrational real roots and four complex roots.

In the table which follows, the first column contains  $n$ , the second, third, and fourth contain the  $a_n$  for the three real roots  $\alpha^{(1)} \sim -1.444 \dots$ ,  $\alpha^{(2)} \sim 0.429 \dots$ ,  $\alpha^{(3)} \sim 1.233$ .

$n$	$\alpha^{(1)}$	$\alpha^{(2)}$	$\alpha^{(3)}$
0	-2	0	1
1	1	2	3
2	1	3	2
3	3	53	2
4	1	5	4
5	86	1	15
6	63	2	4
7	1006	1	1
8	2	1	7
9	1	1	70
10	3	1	1
11	3	91	7
12	2	1	2
13	3	1	1
14	1	1	8
15	1	5	4

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