

A Simplified Version of the Fast Algorithms of Brent and Salamin

By D. J. Newman*

Abstract. We produce more elementary algorithms than those of Brent and Salamin for, respectively, evaluating e^x and π . Although the Gauss arithmetic-geometric process still plays a central role, the elliptic function theory is now unnecessary.

In their remarkable papers, Brent [1] and Salamin [3], respectively, used the theory of elliptic functions to obtain "fast" computations of the function e^x and of the number π . In both cases rather heavy use of elliptic function theory, such as the transformation law of Landen, had to be utilized. Our purpose, in this note, is to give a highly simplified version of their constructions. In our approach, for example, the incomplete elliptic integral is never used.

We begin as they did with the Gauss arithmetic-geometric process, $T(a, b) = ((a + b)/2, \sqrt{ab})$ which maps couples with $a \geq b > 0$ into same. From the inequalities

$$\frac{(a + b)/2 - \sqrt{ab}}{(a + b)/2 + \sqrt{ab}} = \left(\frac{\sqrt{a} - \sqrt{b}}{\sqrt{a} + \sqrt{b}} \right)^2 \leq \left(\frac{a - b}{a + b} \right)^2,$$

and

$$\frac{(a + b)/2}{\sqrt{ab}} \leq \frac{a}{\sqrt{ab}} = \sqrt{\frac{a}{b}},$$

we see that $T^i(a, b)$ goes to its limiting couple (m, m) ($m = m(a, b)$ the so-called arithmetic-geometric mean) with "quadratic" speed. Indeed, $m(a, b)$ is determined to n places for an i of around $\log \log a/b + \log n$. The $\log \log$ from the $\sqrt{a/b}$ inequality expressing the time till the ratio first goes below 2, and the \log from the $((a - b)/(a + b))^2$ inequality expressing the time for the error squaring to do its job.

Next, we recall Gauss' beautiful formula:

$$m(a, b) = \pi / \int_{-\infty}^{\infty} \frac{dx}{\sqrt{(x^2 + a^2)(x^2 + b^2)}},$$

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which follows from the fact that this (complete) elliptic integral is invariant under T . This fact, that namely

$$\int_{-\infty}^{\infty} \frac{dx}{\sqrt{(x^2 + a^2)(x^2 + b^2)}} = \int_{-\infty}^{\infty} \frac{dt}{\sqrt{(t^2 + ((a+b)/2)^2)(t^2 + ab)}},$$

is a simple consequence of the change of variables $t = (x - ab/x)/2$. Namely, we obtain

$$dt = \frac{x^2 + ab}{2x^2} dx, \quad t^2 + \left(\frac{a+b}{2}\right)^2 = \frac{(x^2 + a^2)(x^2 + b^2)}{4x^2},$$

$$t^2 + ab = \frac{(x^2 + ab)^2}{4x^2},$$

$0 < x < \infty$, so that indeed we have

$$\begin{aligned} \int_{-\infty}^{\infty} \frac{dx}{\sqrt{(x^2 + a^2)(x^2 + b^2)}} &= \int_0^{\infty} \frac{2dx}{\sqrt{(x^2 + a^2)(x^2 + b^2)}} \\ &= \int_{-\infty}^{\infty} \frac{dt}{\sqrt{(t^2 + ((a+b)/2)^2)(t^2 + ab)}}. \end{aligned}$$

Accordingly, a repeated use of this invariance gives

$$\begin{aligned} \int_{-\infty}^{\infty} \frac{dx}{\sqrt{(x^2 + a^2)(x^2 + b^2)}} &= \dots = \int_{-\infty}^{\infty} \frac{dx}{\sqrt{(x^2 + m^2)(x^2 + m^2)}} \\ &= \int_{-\infty}^{\infty} \frac{dx}{x^2 + m^2} = \frac{\pi}{m}, \end{aligned}$$

and this is exactly Gauss' formula.

Actually, it is handier for us to work with what we might call the harmonic-geometric mean which can be defined by $h(a, b) = ab/m(a, b)$ or, alternatively, as the limit under repeated applications of S , rather than T , where

$$S(a, b) = (\sqrt{ab}, 2ab/(a+b)).$$

In these terms Gauss' formula reads

$$h(a, b) = \frac{1}{\pi} \int_{-\infty}^{\infty} \frac{dx}{\sqrt{(1 + x^2/a^2)(1 + x^2/b^2)}}.$$

The only place that we actually use this formula is to establish the asymptotic formula:

$$h(N, 1) = \frac{2}{\pi} \log 4N + O(1/N^2).$$

(This simple-looking formula certainly deserves an elementary proof independent of elliptic integrals, but we are unable to find one.)

So begin with

$$h(N, 1) = \frac{2}{\pi} \int_0^{\infty} \frac{dx}{\sqrt{(1 + x^2)(1 + x^2/N^2)}}$$

and observe that the map $x \rightarrow N/x$ leaves the integrand invariant. Thereby, we conclude

$$\int_0^{\sqrt{N}} \frac{dx}{\sqrt{(1+x^2)(1+x^2/N^2)}} = \int_{\sqrt{N}}^{\infty} \frac{dx}{\sqrt{(1+x^2)(1+x^2/N^2)}}$$

which gives us

$$\begin{aligned} h(N, 1) &= \frac{4}{\pi} \int_0^{\sqrt{N}} \frac{dx}{\sqrt{(1+x^2)(1+x^2/N^2)}} \\ &= \frac{4}{\pi} \int_0^{\sqrt{N}} \frac{1}{\sqrt{(1+x^2)}} (1 - x^2/2N^2 + O(x^4/N^4)) dx \\ &= \frac{4}{\pi} \int_0^{\sqrt{N}} \left(\frac{1}{\sqrt{1+x^2}} - \frac{x}{2N^2} \right) dx + O(1/N^2) \\ &= \frac{4}{\pi} (\log(\sqrt{N} + \sqrt{N+1}) - 1/4N) + O(1/N^2) \end{aligned}$$

and so, since

$$\sqrt{N} + \sqrt{N+1} = 2\sqrt{N} \left(1 + 1/(2N + 2\sqrt{N(N+1)}) \right),$$

we obtain

$$\log(\sqrt{N} + \sqrt{N+1}) = \log 2\sqrt{N} + 1/4N + O\left(\frac{1}{N^2}\right),$$

which together with the previous gives

$$h(N, 1) = \frac{4}{\pi} \log 2\sqrt{N} + O\left(\frac{1}{N^2}\right) = \frac{2}{\pi} \log 4N + O\left(\frac{1}{N^2}\right),$$

as required. (This result can also be found in [2].)

Summarizing, then, we have produced a fast method for obtaining n places of $2 \log 4N/\pi$ (if N is of the size c^n). But, and here is the trick, this combination of π and the logarithm can be used to yield both of them separately, and we can thereby rederive both Salamin's and Brent's results.

To obtain π we examine the difference, $h(N+1, 1) - h(N, 1)$, and observe that $N(h(N+1, 1) - h(N, 1)) = 2/\pi + O(1/N)$ which gives n place accuracy for π if we choose, e.g., $N = 2^n$.

For the logarithm, on the other hand, we look to the quotient, $h(N+1, 1)/h(N, 1)$. This time we obtain

$$N \left(\frac{h(N+1, 1)}{h(N, 1)} - 1 \right) = N \frac{\log(1 + 1/N) + O(1/N^2)}{\log 4N + O(1/N)} = \frac{1}{\log 4N + O(1/N)}.$$

From this we will be able to evaluate $\log x$ throughout the interval (3, 9), and so, of course, throughout any interval. And thereby, we will be able to obtain e^x , the inverse function, by the usual use of the (fast) Newton iteration scheme.

To obtain $\log x$, then, in the interval (3, 9), we first calculate $N = \frac{1}{4}x^n$, a process that takes only $\log n$ multiplications. But then the above formula becomes, upon substitution of this value of N ,

$$\frac{1}{4}nx^n \left(\frac{h(\frac{1}{4}x^n + 1, 1)}{h(\frac{1}{4}x^n, 1)} - 1 \right) = \frac{1}{\log x} + O\left(\frac{n}{x^n}\right) = \frac{1}{\log x} + O\left(\frac{n}{3^n}\right)$$

which does give the desired n place evaluation of $\log x$.

This trick of “differencing” $h(N + 1, 1)$ and $h(N, 1)$, of course, carries a price. Thus we must compute these two quantities to $2n$ places and so the running time is around twice as long as the corresponding ones of Brent and Salamin.

Mathematics Department
Temple University
Broad and Montgomery Ave.
Philadelphia, Pennsylvania 19122

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